Verifying ParamGen A Case Study in Scientific Software Abstraction and Modeling

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Improving Scientific Software Conference, April 17, 2023, Boulder, CO

Formal Specification

Describing software using a precise and high-level language.

► Helps design maintainable and reliable software.



Verifying ParamGen

What is ParamGen?

An infrastructure library that generates default input parameter files for ESM components.

- Similar to CIME.nmlgen
- Supports any format (not just nml)
- Supports any template format (not just xml)
- Supports arbitrary Python expressions in template files!



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An infrastructure library that generates default input parameter files for ESM components.

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- Supports any format (not just nml)
- Supports any template format (not just xml)
- Supports arbitrary Python expressions in template files!
- Initially developed for MOM6 in CESM.
- ▶ Used by MOM6, CAM, NUOPC, and WW3 of CESM.



What is ParamGen?





ParamGen Terminology

\downarrow branch

label $ ightarrow$	NIGLOBAL:
guard $ ightarrow$ value $ ightarrow$	<pre>\$OCN_GRID == "g16": 320</pre>

expandable variable ightarrow

\$OCN_GRID



ParamGen is flexible!

- Expandable variables and Python expressions can appear in labels, guards, values.
- ► There can be multiple labels along a branch.
- ► There can be multiple (nested) guards along a branch.
- ► No restriction on the ordering of guards and labels.



ParamGen schema!

Schema:

- There must be at least one label along a branch.
- If a key is a guard, all its siblings must also be guards.

```
NIGLOBAL:
    description:
    "grid points in x-dir."
    $OCN_GRID == "g16":
    320
```

invalid!

```
NIGLOBAL:
  description:
    "grid points in x-dir."
  value:
    $OCN_GRID == "g16":
    320
```

valid



ParamGen reduce() method

def reduce(data):

return data

```
if isinstance(data, dict):
    # (1) Expand vars in keys
    data = expand_key_vars(data)
```

```
# (2) Evaulate guards (if applicable)
if is_guarded_dict(data):
    data = impose_guards(data))
```

```
# (3) Call reduce recursively
else:
   for key in data:
        data[key] = reduce(data[key])
```

else: # (4) expand vars, apply formulas data = expand_vars(data)

► Traverse all branches recursively.

- Apply in-place modifications:
 - expand variables.
 - evaluate formulas.
 - impose guards.



How to ensure reliability?

Sources of complexity

- Recursive algorithm with in-place data modifications.
- Great flexibility in template layout.



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Testing

Routine unit and integration tests.



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Sources of complexity

- Recursive algorithm with in-place data modifications.
- Great flexibility in template layout.

Testing

Routine unit and integration tests.

"Testing can be used to show the presence of bugs, but never to show their absence!."

— Edsger W. Dijkstra



Proving programs correct

Formal methods: mathematically rigorous techniques and tools for the specification, design and verification of software. (Butler, R. W., 2001)



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Formal methods: mathematically rigorous techniques and tools for the specification, design and verification of software. (Butler, R. W., 2001)

Alloy: an open source language and tool for software modeling.

- ► The Alloy language is simple, precise, powerful, and elegant!
- This allows us to model software and algorithms at a high-level. (at the design level, and above the code level)



The Alloy Language in a few slides

Alloy

A software modeling and analysis tool with a **declarative** language that combines first-order logic and relational calculus.

Imperative Programming:

- Describes *how* a program does something step by step.
- Sequential statements change a program's state.

Declarative Programming:

- Describes what a program should accomplish.
- There is no state changes or mutation.



Declarative Programming

- ► There is no state changes or mutation.
- So, how do we represent dynamic behavior?

$$\mathbf{x} = \mathbf{0} \\ \mathbf{x'} = \mathbf{x} + \mathbf{1}$$



Declarative Programming

- ► There is no state changes or mutation.
- So, how do we represent dynamic behavior?

$$\mathbf{x} = 0$$

 $\mathbf{x'} = \mathbf{x} \cdot \mathbf{add} [1]$



In Alloy, everything is a set! Common operators are reserved for set operations:

- + : set union
- : set difference
- = : set equality
- & : set intersection
- in : subset



▶ In alloy we use *signature* declarations to introduce concepts (sets of atoms).

```
sig Value {}
sig Key {
    map: Value
}
```



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► Signature fields, e.g., map, resemble OOP class members.

k.map // returns the value instance k maps to.



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```
sig Value {}
sig Key {
    map: Value
}
```

Signature fields, e.g., map, resemble OOP class members.

k.map // returns the value instance k maps to.

map is a relation:

map : Key -> Value



In Alloy, everything, including sets, is a relation! Relational operators:

- P
 ightarrow Q : arrow product
 - P.Q : dot product
 - $^{\wedge}p$: transitive closure
 - *p : reflexive t.c.

- \sim : transpose
- <: : domain restriction
- :> : range restriction
- ++ : override



Arrow product:

 $P \rightarrow Q$: every combination of tuples (p,q) for $p \in P$ and $q \in Q$.



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Example:

 $P \rightarrow \text{lone } Q$: each member of P maps to zero or one member of Q.



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Relation multiplicities:

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Example:

 $P \rightarrow$ lone Q : each member of P maps to zero or one member of Q.

Multi-arity relations:

 $P \rightarrow Q \rightarrow R$



Modeling ParamGen in Alloy and verifying its correctness

Modeling ParamGen in Alloy

How to model a system in Alloy in an agile and incremental development style:

- 1. Specify the structure (i.e., the signatures)
- 2. Specify dynamic behavior.
- 3. Inspect instances.
- 4. Eliminate design flaws (e.g., by adding more constraints).
- 5. Analyze rigorously.



Main ParamGen Structure: Template

Stored in nested Python Dictionaries.





```
sig Dict {
    contents: set Key
}
abstract sig Key {
    var map : lone Dict+Value
}
sig Value {}
```



```
sig Dict {
   contents: set Key
}
abstract sig Key {
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}
sig Value {}
```

```
sig Label extends Key {}
sig Guard extends Key {}
sig Root extends Key {}
    {no this.~contents}
```



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sig Dict {
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sig Value {}
```

```
sig Label extends Key {}
sig Guard extends Key {}
sig Root extends Key {}
    {no this.~contents}
```





```
// each item can have at most one key
all i: Dict+Value |
    lone i.~map
// each key can be owned by one dict
all k: Key |
    lone k.~contents
```



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all k: Key |
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```
lone k.~contents
```

```
// if a dict key is a guard,
// all dict keys must be so.
all d: Dict |
    {some Guard & d.contents implies
        d.contents in Guard}
```



```
pred invariants {
   // each item can have at most one key
    all i: Dict+Value
       lone i. map
   // each key can be contained by only one dict
    all k: Key
       lone k. contents
   // if a dict key is a guard. all dict keys must be so.
    all d: Dict
        {some Guard & d.contents implies d.contents in Guard}
   // map.*contents relation is acyclic
    no iden & ^(map.*contents)
   // all values must be preceded by a label
    all v: Value
        some v.^(~map.*~contents) & Label
```



```
pred reduce[data: Dict+Value]{
```

```
data in Dict implies {
```

// (1) Expand vars in keys
expand_vars[data.contents]

```
// (2) Evaluate guards
is_guarded_dict[data] implies
impose_guards[data]
```

```
// (3) Call reduce recursively
else
   all key : data.contents |
     key.map' = key.map and
     reduce[key.map']
}
else
   // (4) Expand vars
   expand_vars[data]
```

```
def reduce(data):
```

```
if isinstance(data, dict):
```

```
# (1) Expand vars in keys
data = expand_key_vars(data)
```

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# (2) Evaulate guards (if applicable)
if is_guarded_dict(data):
    data = reduce(impose_guards(data))
```

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# (3) Call reduce recursively
else:
   for key in data:
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else:
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```

return data



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```
// Nondeterministically select one of the
// guards and drop all other branches
pred impose_guards[d: Dict]{
   let pkey = d.~map {
      some g: d.contents {
        pkey.map' = g.map and g.map'=none
        (d.contents-g).map' =
            (d.contents-g).map
        reduce[pkey.map']
   }
}
```

```
def impose guards(self, data dict):
    def eval guard(guard):
        """returns true if a guard evaluates to true."""
        assert isinstance(
            guard, str
        ). "Expression passed to eval guard must be string."
        if has unexpanded var(guard):
            raise RuntimeError("...")
        guard evaluated = eval formula(guard)
        assert isinstance(guard evaluated, bool)
        return guard evaluated
    if not ParamGen.is guarded dict(data dict):
        return data dict
    guards eval true = \begin{bmatrix} 1 \\ \# \end{bmatrix} # list of guards that evaluate to true.
    for guard in data_dict:
        if guard == "else" or _eval_guard(str(guard)) is True:
            guards eval true.append(guard)
    if len(guards_eval_true) > 1 and "else" in guards_eval_true:
        guards eval true.remove("else")
    elif len(guards eval true) == 0:
        return None
```

```
if self._match == "first":
    roturn data_dict[guards_eval_true[0]]
if self._match == "last":
    roturn data_dict[guards_eval_true[-1]]
raiseRund'Unknown match option.")
```



```
let expand_vars[expr] {
    no expr & varsToExpand'
}
```

```
def expand_vars(expr, expand_func):
    if expand_func is None:
        return expr
    assert isinstance(
        expr, str
    ). "Expression passed to expand vars must be string."
    expandable vars = re.findall(r''(\s)_{w+}), expr)
    for word in expandable vars:
        ws = word.strip().
            replace("$", "").replace("{", "").replace("}", "")
        word expanded = expand func(ws)
        accort (
            word expanded is not None
        ). "Cannot determine the value of {},".format(word)
        # enclose with quotes if expanded var is a string...
        if isinstance(word expanded, str) and word[1] != "{":
            word expanded = '"' + word expanded + '"'
        else:
            word_expanded = str(word_expanded)
        expr = re.sub(
            r''(\s)b'' + ws + r''(b)(\s)(f'' + ws + r''))''.
            word expanded, expr.
    return expr
```



Modeling ParamGen Dynamics in Alloy / Abstract Representation

- Abstract representation of lower-level details allows us to focus on high-level software and algorithm design aspects.
 - Quick prototyping and early detection of design flaws!
 - Incremental modeling!



Modeling ParamGen Dynamics in Alloy / Abstract Representation

"The hard part of building software is the specification, design, and testing of conceptual construct, not the labor of representing it and testing the fidelity of the representation."

- FP Brooks. "No silver bullet". (1987)



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"What matters is the fundamental structure of the design. If you get it wrong, there is no amount of bug fixing and refactoring that will produce a reliable, maintainable, and usable system."

- D Jackson. "The essence of software". (2021)



Back to ParamGen Dynamics...





Bounded Model Checking

```
reduce[r] implies {
    invariants and // All invariants still remain true
    r.map' in Dict // The result of postcondition should be a dict
    all ai: r.*(map'.*contents) { // check active columns
        // all labels map to some Values or Dictionaries
        {ai in Label implies ai.map in Value+Null+Dict}
        // no active quard remains
        ai not in Guard
        // all keys should lead to a value
        {ai in Key implies some ai. (map.*contents) & Value}
        // all values should have a label (varname)
        {ai in Value implies some ai. ( map. * contents) & Label}
        // no remaining vars to expand
        ai not in VarsToExpand'}
```



Bounded Model Checking

Executing "Check postconditions for 8 but 2 steps" Solver=glucose 4.1(jni) Steps=1..2 Bitwidth=4 MaxSeq=7 SkolemDepth=1 Symmetry=20 1..2 steps. 144942 vars. 657 primary vars. 338798 clauses. 31625ms. No counterexample found. Assertion may be valid. 97990ms.



Specifying a system helps us understand it.

"It's a good idea to understand a system before building it, so it's a good idea to write a specification of a system before implementing it."

- L Lamport. "Specifying Systems". (2002)

"Code is a poor medium for exploring abstractions."

— D Jackson. "Software Abstractions". (2006)



Specifying Scientific Software

- Specifying a system beforehand is not always possible in computational science.
 - Sometimes, we are handed over pre-developed code for integration.
 - Sometimes, complexity is underestimated and/or targets change.
 - Often, we work with legacy code.



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- Specification and modeling may be useful at later stages of development as well.
 - Modeling ParamGen in Alloy helped me better understand it (even though I was the original developer).



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 - Sometimes, complexity is underestimated and/or targets change.
 - Often, we work with legacy code.
- Specification and modeling may be useful at later stages of development as well.
 - Modeling ParamGen in Alloy helped me better understand it (even though I was the original developer).
- After developing the Alloy model of ParamGen:
 - ▶ I refactored the **reduce()** method and halved the LOC.
 - I have identified flaws in template handling and sanity checks.



Summary

Formal specification and software modeling can help us develop:

- Reliable software.
 - Testing is crucial, but incomplete.
 - Formal methods can be complementary to testing.
- Maintainable and efficient software.
 - Scientists and engineers are accustomed to working with models anyway.
 - With automatic, push-button analysis, one can focus on modeling and design aspects.
 - ▶ Design aspects matter. More so than lower level implementation details.



Announcement

November 9-10, 2023 Workshop on Correctness and Reproducibility for Climate and Weather Software

website: ncar.github.io/correctness-workshop/

Scope: Testing, Statistical Approaches, Formal Methods, Verification and Validation... Venue: NCAR Mesa Lab & Virtual Thanks! altuntas@ucar.edu